Distributed Shared Memory

CS239
Advanced Operating Systems

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Lecture Source Materials

Singhal, Ch. 10: **Distributed Shared Memory**

Towards Transparent and Efficient Software DSM, Scales et al, SOSP 97.

Cashmere-2L:Software Coherent Shared Memory on a Clustered Remote-Write Network, Stets et al, SOSP 97.

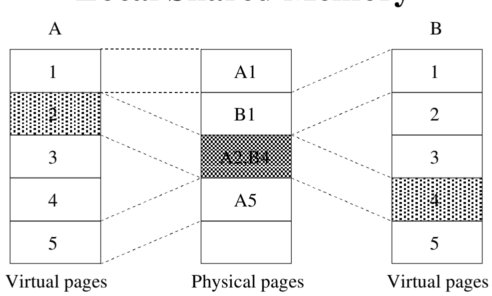
Outline

- Introduction and Motivation
- Basic Approaches
- Memory Coherence and Protocols
- Design Issues
- Example Systems
- Recent Results in DSM Research
- Summary

Introduction and Motivation

- Local shared memory
- Basis for distributed computation?

Local Shared Memory



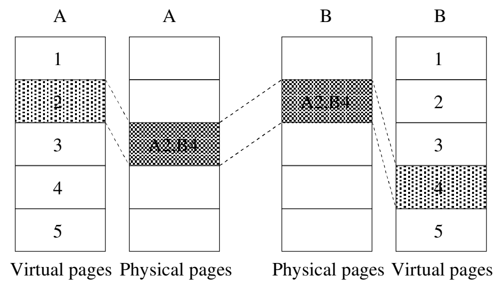
Local Shared Memory

- features
  - "cheap" IPC for cooperating processes
  - synchronization left to processess
  - useful in tightly-coupled multiprocessors
    - exploit concurrent execution
    - exploit common physical memory pool

Local Shared Memory, cont.

- potential problems
  - coherence
    - read/write access serialized by memory hardware
  - VM page thrashing
    - shared pages are “stickier” than non-shared pages

Distributed Shared Memory



Motivation

- DSM hopes to use loosely-coupled processors to get benefits of SM without paying high $$ cost for multiprocessor

DSM Advantages (claimed)

- simpler abstraction than message passing
- easier to access complex data structures
- exploits locality of reference properties
- pools all physical memory in one huge memory
- scalable (e.g., no common bus

DSM disadvantages (possible)

- Applications know a lot about synchronization methods/tools
- Performance will suffer if too much remote memory access

The Challenge

Can cooperating processes on processors that do not share a common physical memory be presented a coherent shared memory space?

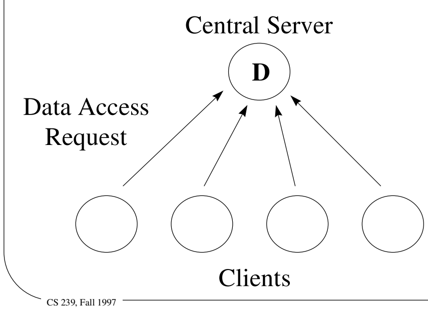
Yes,

… and sometimes with a reasonable cost in performance!

Approaches to DSM

- Central server
- Migration
- Read replication
- Full replication

Central Server DSM



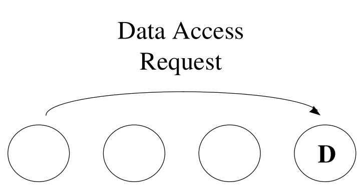
Central Server DSM

- A central server maintains all data pages
  - returns data on reads
  - makes updates on writes
- Care needed to avoid duplicating writes
- Server duties can be split between several sites, by statically partitioning the shared memory page pool

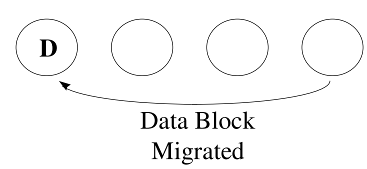
Pros/Cons of Central Server DSM

+ Relatively simple to implement
+ Synchronization methods pretty clear
– Most data references are remote
– Prone to bottlenecks

Migration DSM



Migration DSM



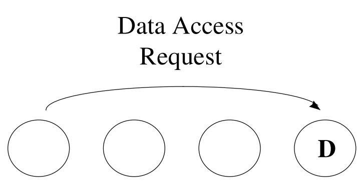
Migration DSM

- Data lives where most recently used
- When “off-site” memory is requested, it migrates to the requesting site
- Granularity is usually block or page level
- With locality of reference, works well
- But possibility of thrashing
- Servers keep track of data locations

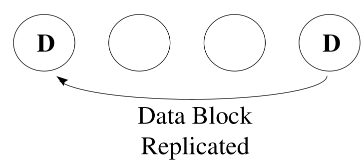
Pros/Cons of Migration DSM

+ If locality good, most accesses can be local
+ Synchronization behavior fairly easy
– Between-node thrashing dangers
– Possible bottlenecks at location servers
– May put heavy burden on communications

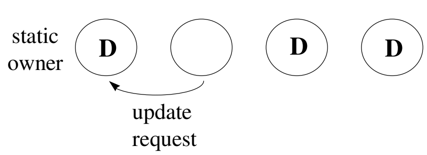
Read Replication DSM



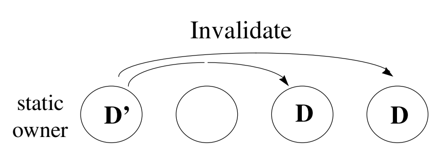
Read Replication DSM



Read Replication DSM with update



Read Replication DSM with update



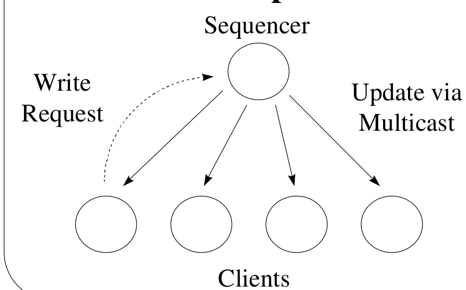
Read-Replication DSM

- Multiple copies of readable blocks stored
- If block is written, invalidate other copies
- Each block has owner that permanently stores it (and it is the only updatable replica)
- If writes are frequent, invalidation can

Pros/Cons of Read Replication DSM

+ Reads are usually local
+ Synchronization not too difficult
– Requires invalidations
– Possibly heavy burden on owners of popular pages
– Writes are only local for page owner

Full Replication DSM with update



Full-Replication DSM

- Full-service replicas of data blocks/pages stored at several machines
- Requires consistency mechanism
  - for same reasons as replicated files
  - different mechanisms may be appropriate

Pros/Cons of Full Replication DSM

+ Access usually local
+ Bottlenecks generally less likely
– Synchronization tricky
– Requires full replication mechanism
– Likely to be overkill unless data sharing very common

Memory Coherence in Replicated DSM Schemes

If more than one copy of a data item exists, problems arise when data is written and "later" read

Consistency model describes what level of consistency a DSM system provides

Consistency Models for DSM

- Strict consistency (as if only one copy)
- Sequential consistency (serializability)
- General consistency (eventual consistency)
- Processor consistency (per machine)
- Weak consistency (up to programmer)

Protocols to Provide DSM Coherence

- Write-invalidate protocol
  - Invalidate other copies before write
  - Expensive for those other copies
- Write-update protocol
  - Update other copies on write
  - Higher network traffic results
  - Requires some kind of locking

Controlling Thrashing in DSM

How to prevent 2 sites sharing a page for write from thrashing the net?

One way is to lock a written page down at a site for some period Δ

If another site wants to write it, that site must wait for Δ to expire

Used in Mirage system

Design Issues

- Granularity of shared item
- Page replacement policy
- Exploiting semantic hints

Granularity in DSM

How big should the unit of shared memory be?

- Should be related to hardware's page size
- Size of page has many important implications...

Pros/Cons of Large Pages for DSM

\+ Takes advantage of locality of reference

\+ Minimizes cost of running paging code

\+ Can minimize communications cost

- If lots of data being moved

– Leads to greater contention

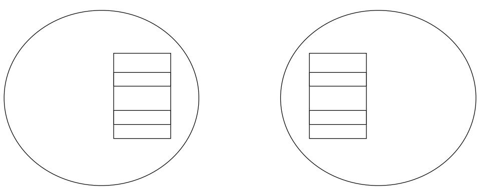
Pros/Cons of Small Page Size for DSM

+ Less contention
+ Minimizes *false sharing*
+ Lower cost to move small data objects
– Greater overhead per byte

False Sharing in DSM

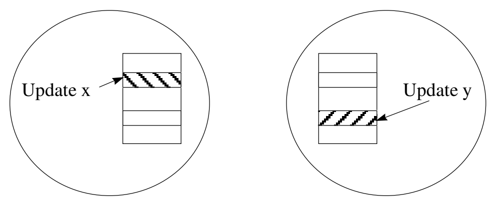
- Occurs when an apparent, but not actual, write conflict occurs between two copies of a page
- Caused by viewing replication unit as a page, rather than a memory location (or object)

False Sharing Example



Site A Site B

False Sharing Example, Con't



Site A Site B

There appears to be a conflict, but there really isn't

Effects of False Sharing

- Depends on approach used
  - In migration systems, causes more page migration
  - In write-invalidate replication systems, causes more invalidations
  - In update-propagation replication systems, causes more locking

Page Replacement in DSM

- When a local node needs to bump a shared memory page, what does it do?
- LRU alone not a sufficient answer
- Must consider private vs. shared pages
  - And replication issues
  - And read-only pages
- LRU within classes often used

Exploiting Semantic Hints

- Write-once
- Private
- Write-many
- Result
- Synchronization
- Migratory
- Producer-consumer

Semantic hints, cont.

- Read-mostly
- General read-write
  - invalid
  - unowned
  - owned exclusively
  - owned nonexclusively

Some DSM Systems

- Ivy (Yale)
- Mirage (UCLA)
- Clouds (Georgia Tech)
- Munin (Rice U.)
- Memnet (U. of Delaware)
- Telegraphos (Greece)

Recent DSM Research

- SOSP 97
  - Shasta
  - Cashmere-2L
- Emergence of low-cost small-scale SMP with hardware shared memory
- Low-latency remote-write networks (e.g., DEC's Memory Channel)

DEC's Memory Channel

- Low-latency hardware interconnect
- Knows about process-process connections (up to 64K)
- Appears as both device (outgoing) and physical memory (incoming)
- Write-only paradigm…
  - how to do a remote read?

Shasta

- Goals
  - transparent execution of standard binaries (no special knowledge of software DSM services)
  - fine-grained sharing (smaller than page-level)

Shasta techniques

- Cache coherence protocol
  - private v shared virtual address space
  - shared states
    - invalid vs shared vs exclusive
    - shared miss
- Code insertion to detect shared miss load/store before it actually occurs

Shasta Techniques, continued

- Avoid OS involvement where possible
  - use inline code for shared misses, not OS/VM interrupt handling
  - use polling for Memory Channel, not standard device interrupt

Cashmere

- Goal
  - two-level shared memory support
    - hardware within SMP "node"
    - software across nodes
  - uses release consistency, a la earlier DSM implementations

DSM Summary

- Basic idea of DSM is…
- Primary problems of DSM are…
- Recent HW trend responses are...